# CSc 461/561 Multimedia Systems Review on TCP/IP Networking

Jianping Pan Spring 2006

2/7/06 CSc 461/561

## Application-oriented view

#### • Applications

- remote login: e.g., telnet
- file transfer: e.g., ftp
- electronic mail: e.g., email
- world-wide web: the Web!

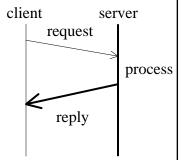
#### • Requirements

- move data from one location to another
- elastic, error-free, in-sequence

### Client-server applications

- E.g., HTTP
  - HTTP client (browser)
    - **GET** /index.html HTTP/1.1
    - Host: www.example.com
    - (parameters)
  - HTTP server (Web server)
    - HTTP/1.1 200 OK
    - (metadata)
    - (data)

2/7/06



3

### Protocols to support

CSc 461/561

- TCP/IP
  - the Internet Protocol Suite
- TCP offers
  - connection oriented
  - reliable, in-sequence, stream-like data transfer
- IP offers
  - addressing and routing; connectionless
  - IP packets may be lost, corrupted, duplicated, reordered

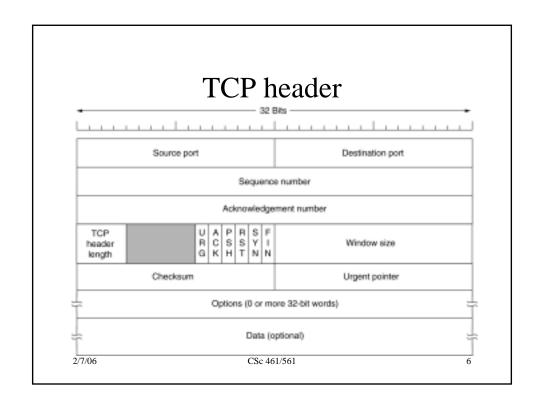
2/7/06 CSc 461/561

link

Ethernet, etc

#### **TCP**

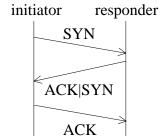
- Connection management
  - through packet handshake (SYN, FIN, ACK)
  - Multiplexing: port number
- Flow, error, congestion control
  - sequence number
  - acknowledgment number
  - window size
  - checksum

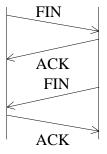


#### Packet handshake

connection establishment

connection release



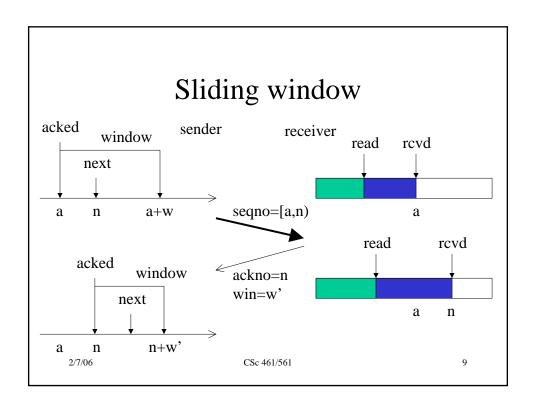


7

2/7/06 CSc 461/561

#### Flow control

- Purpose: pace sender and receiver
  - according to their buffer size
- Receiver's (advertised) window
  - available buffer space
- Sender's window
  - sliding window-based flow control
  - only send data within the window
  - sender's window < receiver's window



#### Error control

- Error detection
  - receiver: sequence number, TCP checksum
  - sender: timeout
- Error notification
  - receiver=>sender: duplicate acknowledgment
- Error recovery
  - sender: end-to-end retransmission

#### Go-back-N retransmission

- Sender: send packet 1, 2, 3, 4, 5, 6
- Receiver: receive packet 1, 2, 3, 5, 6
  - cumulatively acknowledge up to packet 3
- Sender: send packet 7, 8, 9
  - timeout for packet 4; retransmit packet 4
- Receiver
  - cumulatively acknowledge up to packet 9
- Sender: send packet 10, 11, 12, ...

11

12

### Congestion control

- Was not there when TCP was first designed
- Added to TCP since late 80s
  - heavily coupled with flow/error control
  - heavily explored research topics in decades!
- Purpose: pace sender and network
  - competing flows
  - overload network (e.g., output queue)
  - cause packet losses when queues overflow

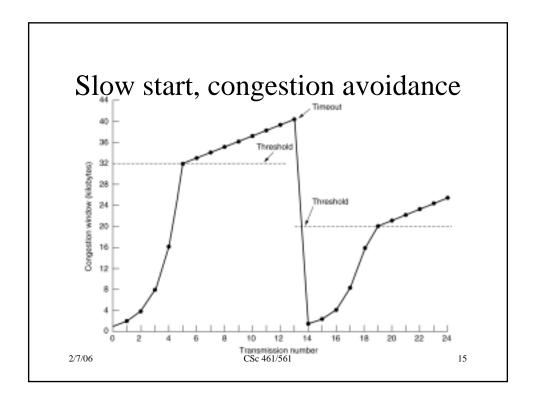
### Congestion window

- Slow start
  - start with an initial congestion window (cwnd)
  - usually initial cwnd = 1 full-size packet
  - double window size every round-trip time until cwnd is above slow start threshold (ssthresh)
- Congestion avoidance
  - increase window size linearly

2/7/06 CSc 461/561 1

### Congestion window: more

- · Back-off
  - when timeout occurs, assume packet loss happened [error control]
  - ssthresh = 0.5 \* current\_cwnd
  - restart with the initial cwnd
  - retransmit with doubled timer [error control]



## Congestion control: more

- Fast retransmit
  - retransmit with 3 duplicate acknowledgments
  - slow start threshold (ssthresh) to be a half of current congestion window (cwnd)
  - restart with the initial congestion window
- Fast recovery
  - similar to Fast retransmit
  - but restart with cwnd = ssthresh

#### **UDP**

- Why TCP is not enough?
  - sometimes TCP is an overkill
  - e.g., loss-tolerant, delay-sensitive applications
- UDP offers
  - connectionless, datagram-like data transfer
  - no reliability, in-sequence guarantee
  - flexibility to plug-in flow/error/congestion etc in application layers

2/7/06 CSc 461/561 17

#### **UDP** header



### This lecture

- A quick review on
  - Internet Protocol Suite
  - TCP
  - UDP
- Bonus question
  - Why TCP header has no "TCP length" field such as "UDP length" in UDP header?

2/7/06 CSc 461/561 19

### Next lecture

- IP
- Why multimedia networking is different?