## 1. Circle **True** or **False** and justify your answer. **No marks will be given unless** there is a correct justification.

- (a) [5 marks] The language  $L = \{w : w \text{ contains } aab \text{ and } aba\}$  defined over  $\Sigma = \{a, b\}$  is regular. True False
- (b) [5 marks] If  $L_1$  and  $L_2$  are regular languages, then so is  $L_1 L_2$ . True False
- (c) [5 marks] Every subset of a regular language is regular. True False
- (d) [5 marks] For every language L over Σ = { a, b }, it is possible to find some regular language that is a subset of L.
  True False
- 2 [20 marks] Given that  $M_1 = (K_1, \Sigma, \Delta_1, s_1, F_1)$  is a FA that accepts a language  $L_1$  and  $M_2 = (K_2, \Sigma, \Delta_2, s_2, F_2)$  is a FA that accepts a language  $L_2$ , describe how to construct a finite automata  $M = (K, \Sigma, \Delta, s, F)$  that accepts  $L_1 \circ L_2$ .

It is fine to draw a visual aid to help in determining your solution. However for full marks, you must precisely describe how to create M.

- 3.(a) [10 marks] State precisely the pumping lemma for regular languages.
- (b) [10 marks] Describe all ways of factoring  $w = a^r b^s c^{r+s}$  as x y z where y is not equal to the empty string.
- 4. [20 marks] Apply the pumping lemma to  $w = a^r b^s c^{r+s}$ . to prove that  $L = \{a^n b^m c^p : n + m \le p\}$  is not accepted by a DFA with k = 2(r+s) states.

5. [20 marks] Use the construction described in class (which is the same as the one in the text) to convert this NDFA to an equivalent DFA:



State	Symbol		Next state

Start state:

Final states:

A picture of your final DFA:  $M = (K, \Sigma, \Delta, s, F)$ .