

Prove the following languages over  $\Sigma=\{0,1\}$  are regular by giving regular expressions for them:

1.  $\{w \text{ contains two or more } 0\text{'s}\}$
2.  $\{|w| = 3k \text{ for some integer } k\}$
3.  $\{w \text{ has } 001 \text{ or } 1010 \text{ as a prefix}\}$
4.  $\{w \text{ contains both } 00 \text{ and } 11 \text{ as substrings}\}$
5.  $\{w \text{ contains } 01001 \text{ and } 0101 \text{ as substrings}\}$
6.  $\{w : w \text{ has at most two consecutive } 1\text{'s}\}$

# Deterministic Finite Automata

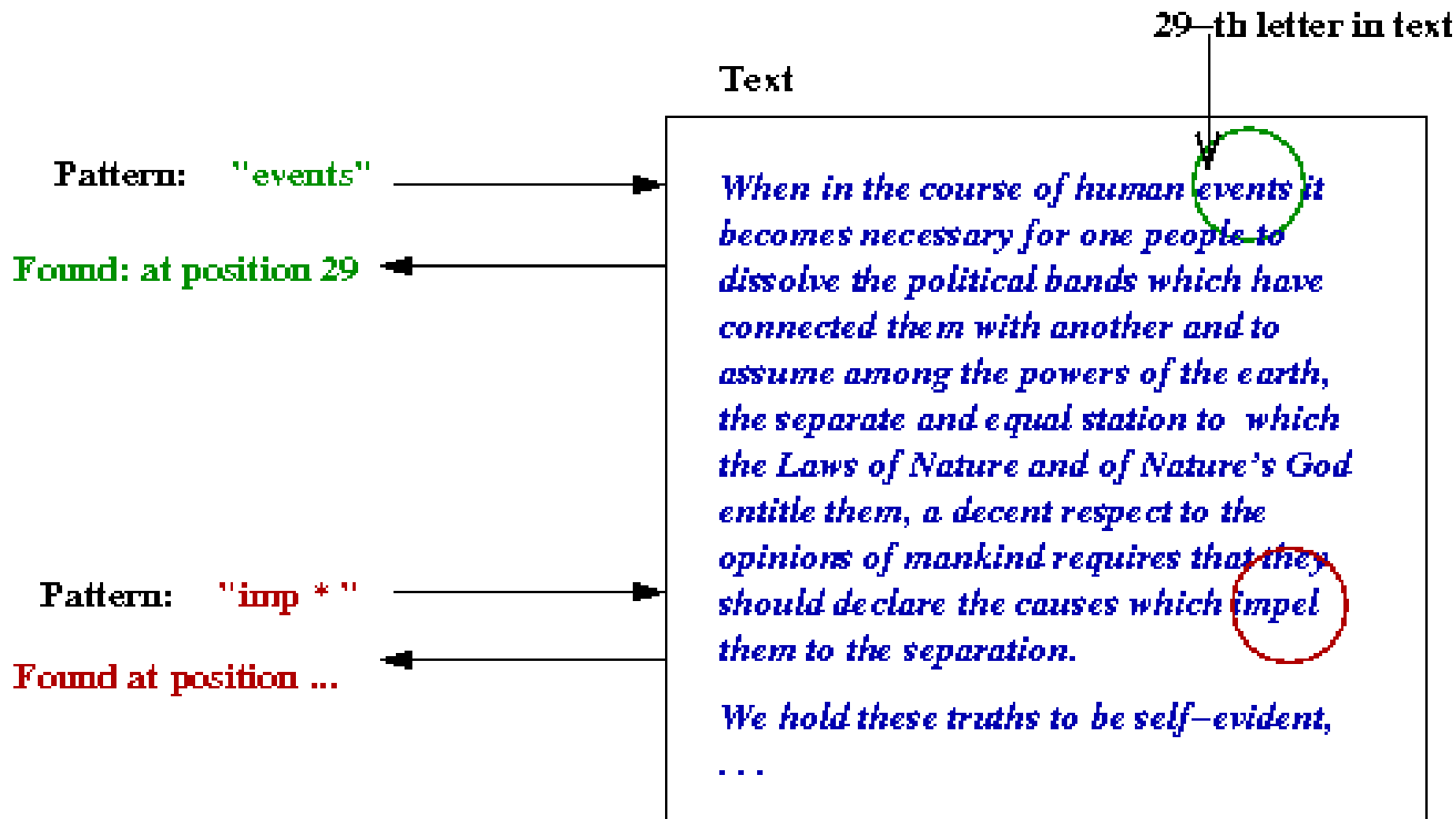
Given a language  $L$ , the **language recognition problem** is:

Given an input string  $w$ , is  $w$  in  $L$ ?

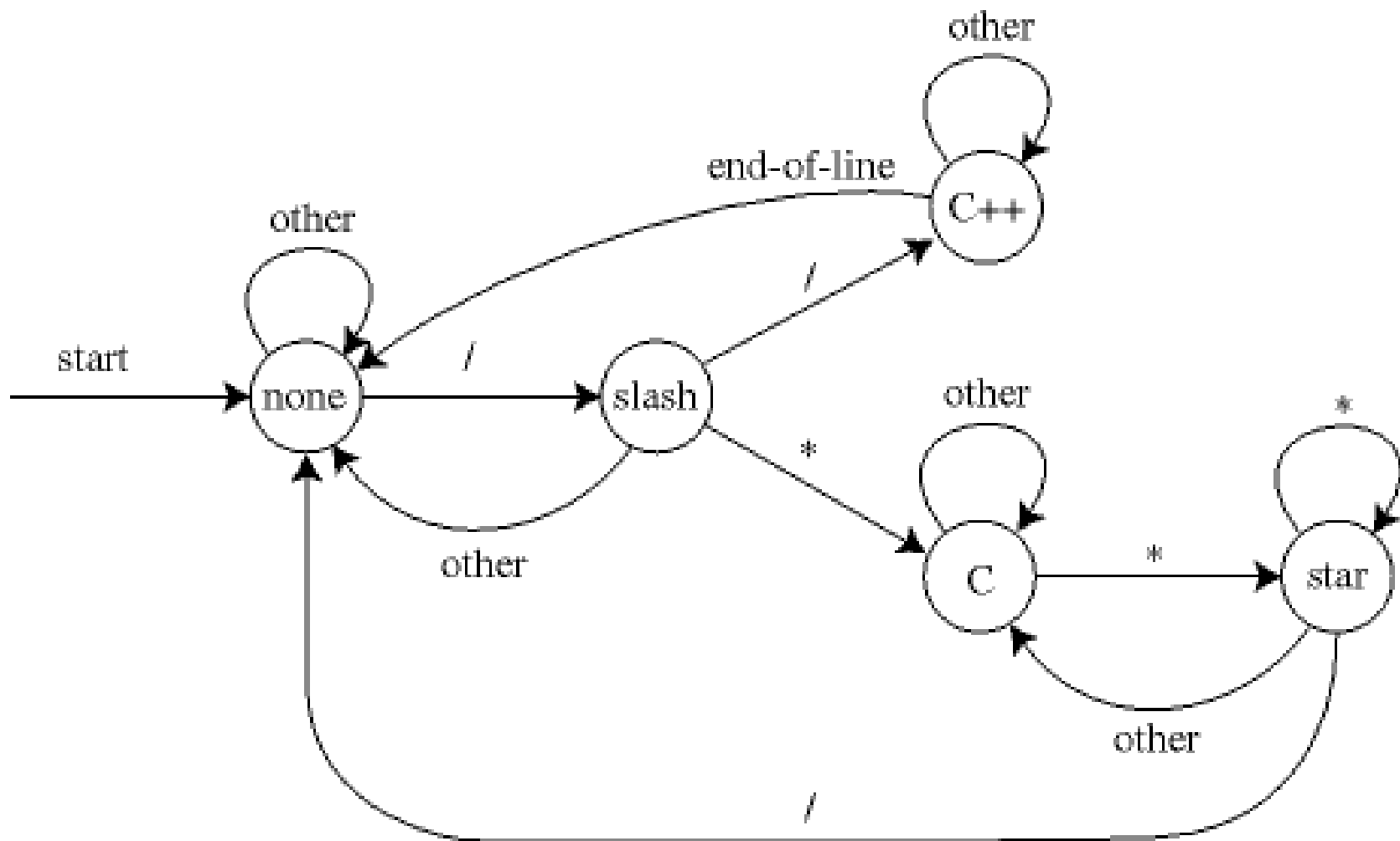
Answer: yes or no.

DFA's provide linear time language recognition algorithms for regular languages (languages which can be defined by a regular expression).

This class introduces DFA's, describes how they work, and defines the mathematical notation used to talk about them.



# To strip comments from a program:



# Applications of DFA's or DFA-like code:

[[www.fbeedle.com/formallanguage/ch04.pdf](http://www.fbeedle.com/formallanguage/ch04.pdf)]

**Compilers and other language-processing systems-** to divide an input program into tokens like identifiers, constants, and keywords and to remove comments and white space.

**Typed commands-** the command language can be quite complex, almost like a little programming language.

**Speech-processing and other signal-processing systems-** to transform an incoming signal.

# More Applications

[[www.fbeedle.com/formallanguage/ch04.pdf](http://www.fbeedle.com/formallanguage/ch04.pdf)]

**Text processors-** to search a text file for strings that match a given pattern. This includes most versions of Unix tools like awk, egrep, and Procmail, along with a number of platform-independent systems such as MySQL.

**Controllers-** to track the current state of a wide variety of finite-state systems, from industrial processes to video games, implemented in hardware or in software.

$L = \{w \in \{a,b\}^* : \text{the number of } a\text{'s in } w \text{ is not divisible by } 3\}.$

1. Give a regular expression for  $L$ .

2. Design a DFA that accepts  $L$ .

Give regular expressions for:

$L_1 = \{w \in \{a,b\}^* : w = a^{2k} b^{2r+1} \text{ for some integers } k, r \geq 1\}$

$L_2 = \{w \in \{a,b\}^* : w \text{ has at least two } a\text{'s and at least two } b\text{'s}\}$



A **Deterministic Finite Automaton** (DFA)

$M$  is defined to be a quintuple

$(K, \Sigma, \delta, s, F)$  where

$K$  is a finite set of **states**,

$\Sigma$  is an alphabet,

$\delta$  is a **function** from  $K \times \Sigma$  to  $K$ ,

$s \in K$  is the **start state**, and

$F \subseteq K$  is the set of **final states**.

Some examples:

1.  $\{w: w \text{ has odd length}\}$
2.  $\{w: w \text{ contains } 0011\}$
3.  $\{w: w \text{ does not contain } 01\}$
4.  $\{w: w \text{ starts and ends with the same symbol } (|w| \geq 1)\}$

$M = (K, \Sigma, \delta, s, F)$ :

A **configuration** of a  $M$  is an element of  $K \times \Sigma^*$ .

For  $\sigma \in \Sigma$ , configuration  $(q, \sigma w) \vdash (r, w)$

$[(q, \sigma w) \text{ yields in one step } (r, w)]$

if  $\delta(q, \sigma) = r$ .

The notation  $\vdash^*$  means yields in zero or more steps.

$M = (K, \Sigma, \delta, s, F)$ :

DFA  $M$  **accepts** input  $w$  if and only if  
 $(s, w) \vdash^* (f, \varepsilon)$  for some  $f \in F$ .

**$L(M)$** , the **language accepted by  $M$**  is  
 $\{ w \in \Sigma^* : (s, w) \vdash^* (f, \varepsilon) \text{ for some } f \in F \}$ .

What language does this DFA accept?

