

CS 330 Lecture 15

- More operational Semantics

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Control Flow

IFTRUE:

$\langle e_1, g, f, p \rangle \Rightarrow \langle v_1, g', f, p' \rangle \quad v_1 \neq 0 \quad \langle e_2, g', f, p' \rangle \Rightarrow \langle v_2, g'', f, p'' \rangle$

$\langle \text{IF}(e_1, e_2, e_3), g, f, p \rangle \Rightarrow \langle v_2, g'', f, p'' \rangle$

IFFALSE:

$\langle e_1, g, f, p \rangle \Rightarrow \langle v_1, g', f, p' \rangle \quad v_1 = 0 \quad \langle e_3, g', f, p' \rangle \Rightarrow \langle v_3, g'', f, p'' \rangle$

$\langle \text{IF}(e_1, e_2, e_3), g, f, p \rangle \Rightarrow \langle v_3, g'', f, p'' \rangle$

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Loops

WHILE-ITERATE

$\langle e_1, g, f, p \rangle \Rightarrow \langle v_1, g', f, p' \rangle \quad v_1 \neq 0$
 $\langle e_2, g', f, p' \rangle \Rightarrow \langle v_2, g'', f, p'' \rangle \quad \langle \text{WHILE}(e_1, e_2, g'', f, p') \rangle \Rightarrow \langle v_3, g''', f, p''' \rangle$

$\langle \text{WHILE}(e_1, e_2, g, f, p) \rangle \Rightarrow \langle v_3, g''', f, p''' \rangle$

WHILEEND

$\langle e_1, g, f, p \rangle \Rightarrow \langle v_1, g', f, p' \rangle \quad v_1 = 0$

e_2 evaluated only for
side effects

$\langle \text{WHILE}(e_1, e_2, g, f, p) \rangle \Rightarrow \langle 0, g', f, p' \rangle$

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Sequential Execution

order of expressions
matters
order of premises
doesn't

$\langle \text{BEGIN}(), g, f, p \rangle \Rightarrow \langle 0, g, f, p \rangle$

$\langle e_1, g_0, f, p_0 \rangle \Rightarrow \langle v_1, g_1, f, p_1 \rangle$

$\langle e_2, g_1, f, p_1 \rangle \Rightarrow \langle v_2, g_2, f, p_2 \rangle$

....

....

$\langle e_n, g_{n-1}, f, p_{n-1} \rangle \Rightarrow \langle v_n, g_n, f, p_n \rangle$

$\langle \text{BEGIN}(e_1, e_2, \dots, e_n), g_0, f, p_0 \rangle \Rightarrow \langle v_n, g_n, f, p_n \rangle$

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Function Application

$f(\text{foo}) = \text{USER}(<\text{x}_1, \dots, \text{x}_n>, e)$

$\text{x}_1, \dots, \text{x}_n$ all distinct

$<e_1, g_0, f, p_0> \Rightarrow <v_1, g_1, f, p_1>$

...

$<e_n, g_{n-1}, f, p_{n-1}> \Rightarrow <v_n, g_n, f, p_n>$

$<e, g_n, f, \{x_1 \rightarrow v_1, \dots, x_n \rightarrow v_n\}> \Rightarrow <v, g', f, p>$

----- APPLY USER

$<\text{APPLY}(\text{foo}, e_1, \dots, e_n), g_0, f, p_0> \Rightarrow <v, g', f, p_n>$

behavior of function doesn't depend on function name, only definition body of a function can't get the formal parameters of its caller

functions assigns to formal parameters changes are not visible outside

p ' can be thrown away after evaluation => VERY IMPORTANT

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Primitive Functions

$f(\text{foo}) = \text{PRIMITIVE}(+)$

$<e_1, g_0, f, p_0> \Rightarrow <v_1, g_1, f, p_1>$

$<e_2, g_1, f, p_1> \Rightarrow <v_2, g_2, f, p_2>$

 $<\text{APPLY}(\text{foo}, e_1, e_2), g_0, f, p_0> \Rightarrow <v_1 + v_2, g_2, f, p_2>$

----- APPLYADD

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Top-level items

$<e, g, f, \{ \}> \Rightarrow <v, g', f, p>$

----- EVALEXP

$<\text{EXP}(e), g, f> \rightarrow <g', f>$

$<e, g, f, \{ \}> \Rightarrow <v, g', f, p>$

----- DEFINEGLOBAL

$<\text{VAL}(x, e), g, f> \rightarrow <g'\{x \rightarrow v\}, f>$

x_1, x_2, \dots, x_n all distinct

----- DEFINE FUNCTION

$<\text{DEFINE}(\text{foo}, <\text{x}_1, \dots, \text{x}_n>, e), g, f> \rightarrow \{g, f \rightarrow \text{USER}(<\text{x}_1, \dots, \text{x}_n>, e)\}$

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