

Minimum weight spanning tree:

Put a non-negative weight on each edge.

Weight of a tree: sum of weights of its edges.

Problem: find a spanning tree of graph  $G$  with minimum weight.

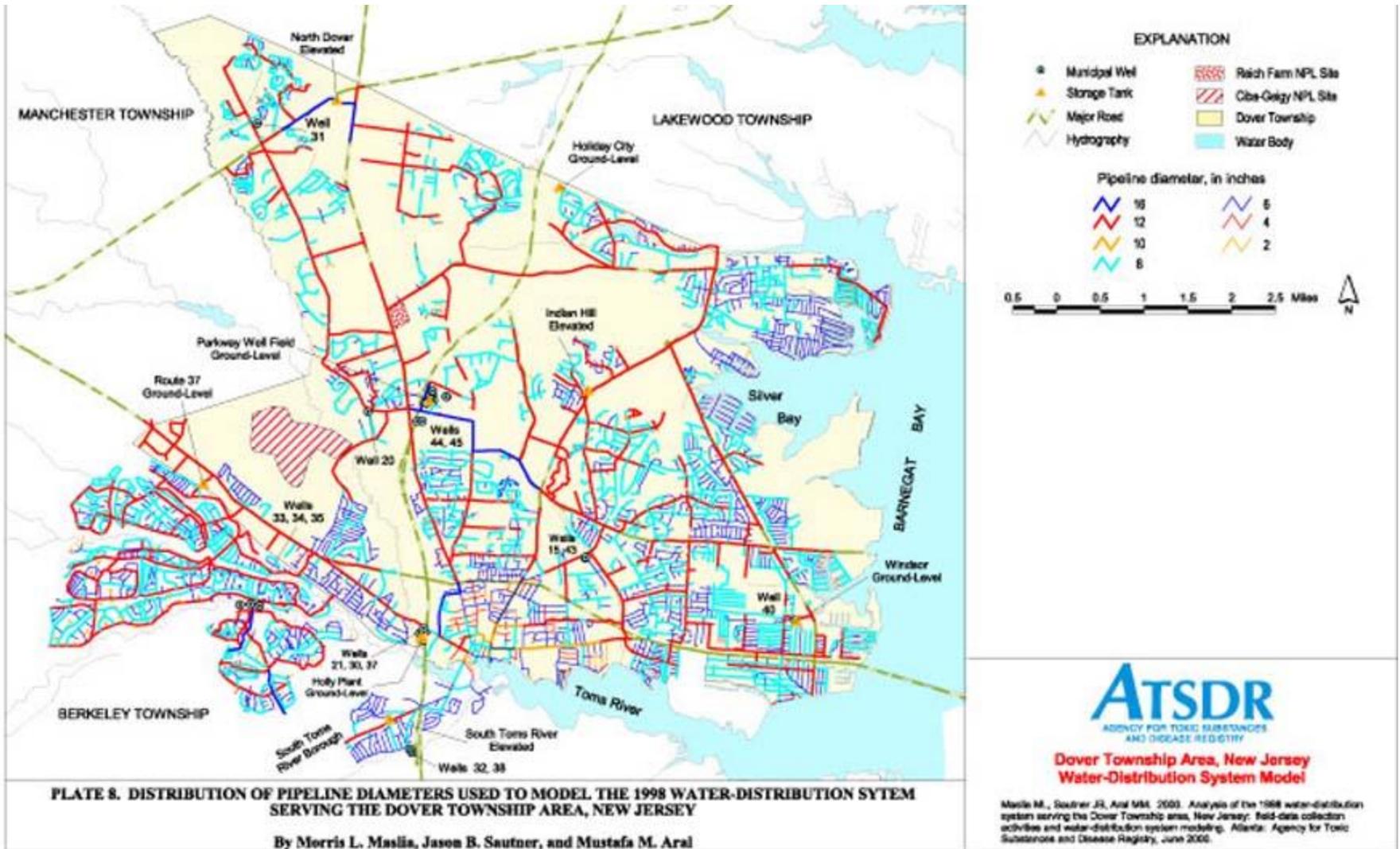
Kruskal's algorithm: sort edges by weight.

Then for each edge: add it to the tree if its endpoints are in different components.

One application:

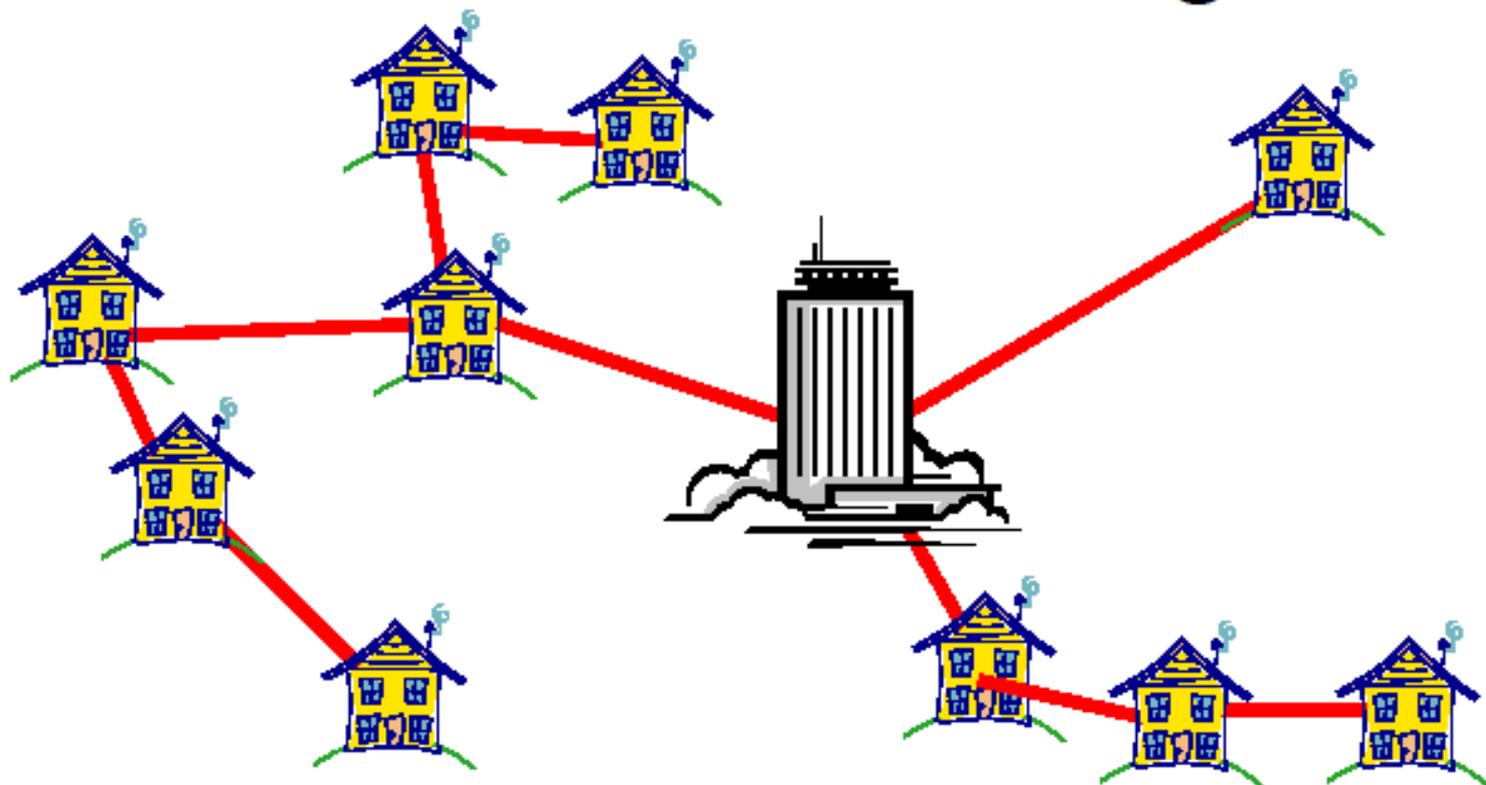
A cable company must install cable to a new neighbourhood. The cables are constrained to be buried along certain paths. The cost varies for different paths. A minimum weight spanning tree gives the cheapest way to connect everyone to cable.

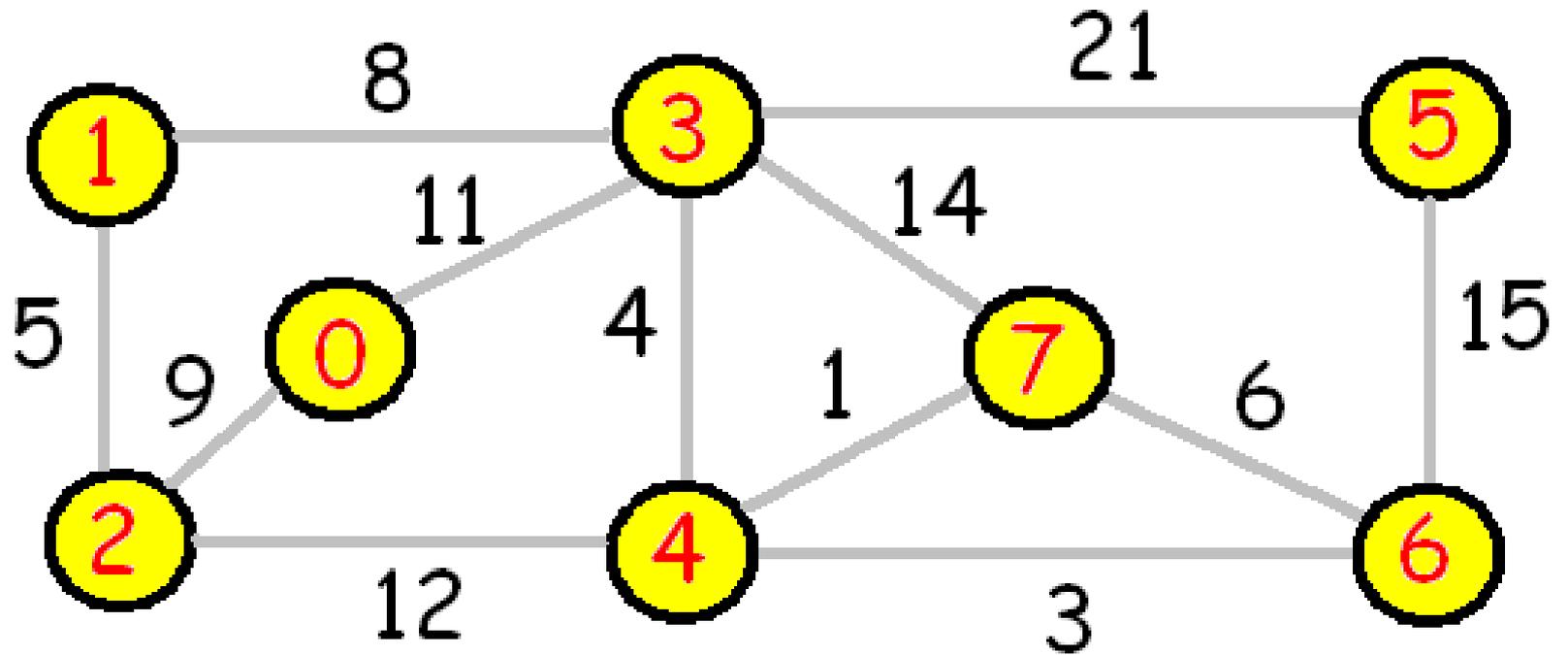
# Water distribution network





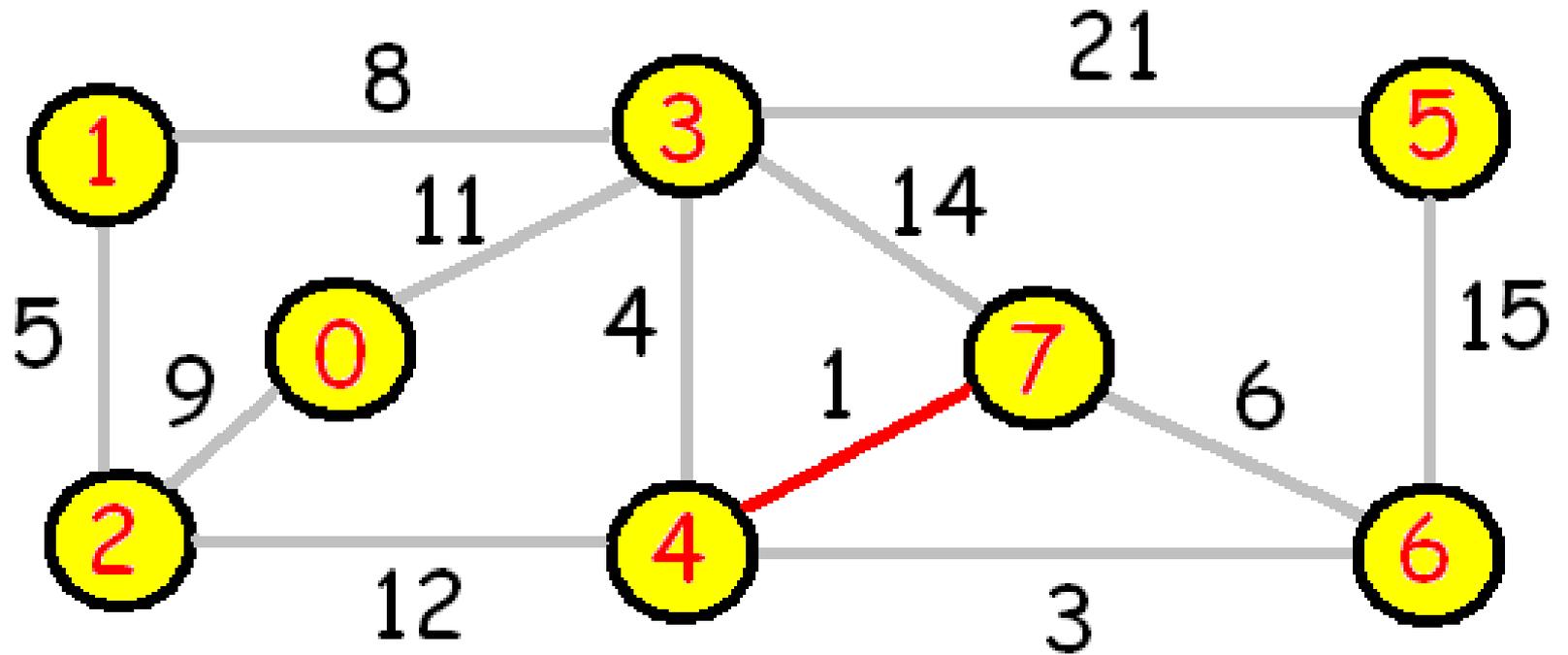
# Minimize Wiring





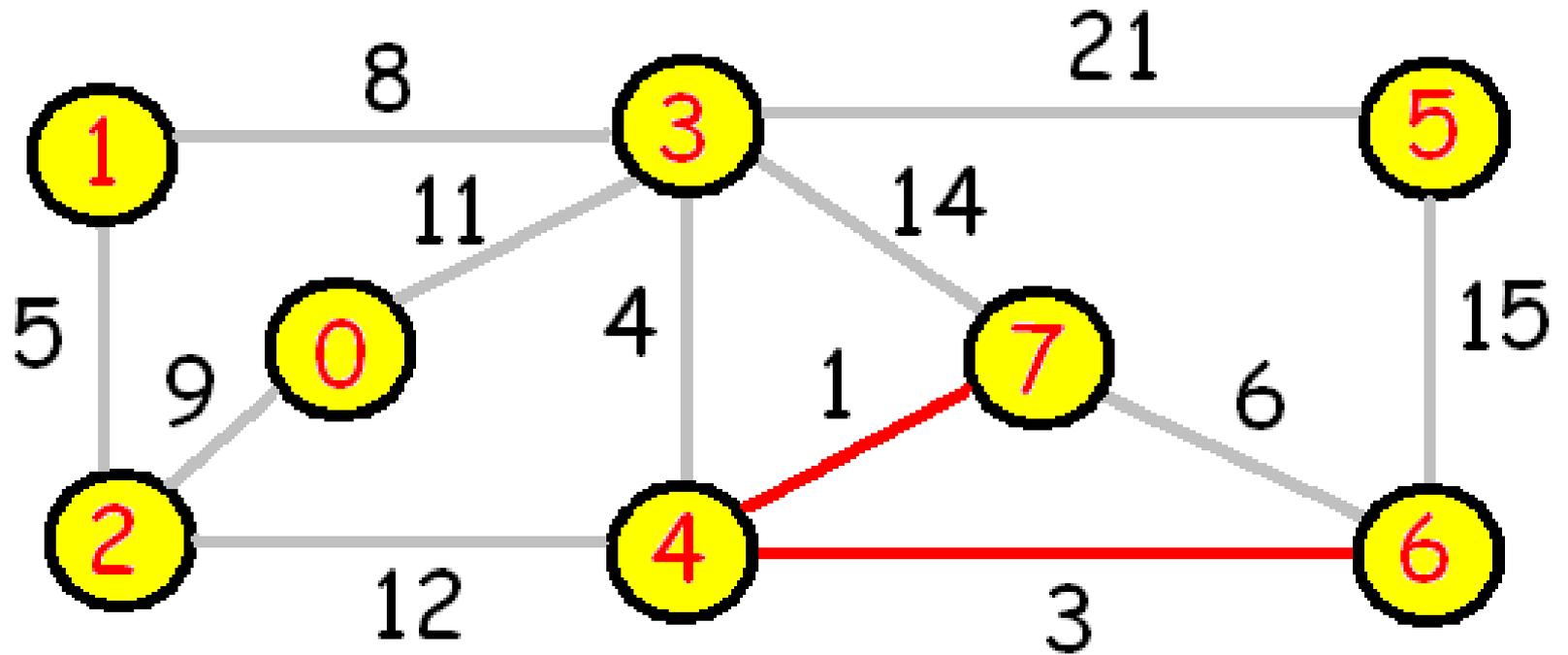
Edge weights:

1, 3, 4, 5, 6, 8, 9, 11, 12, 14, 15, 21



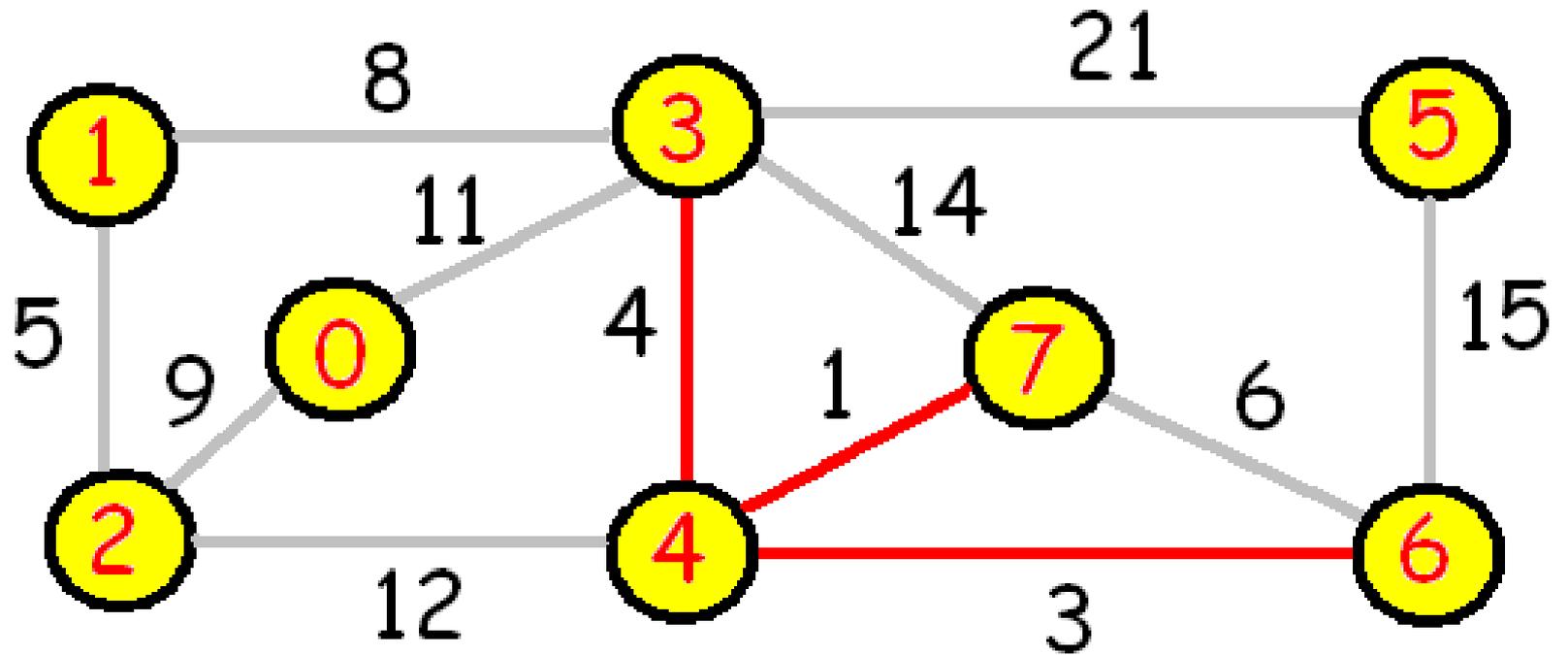
Edge weights:

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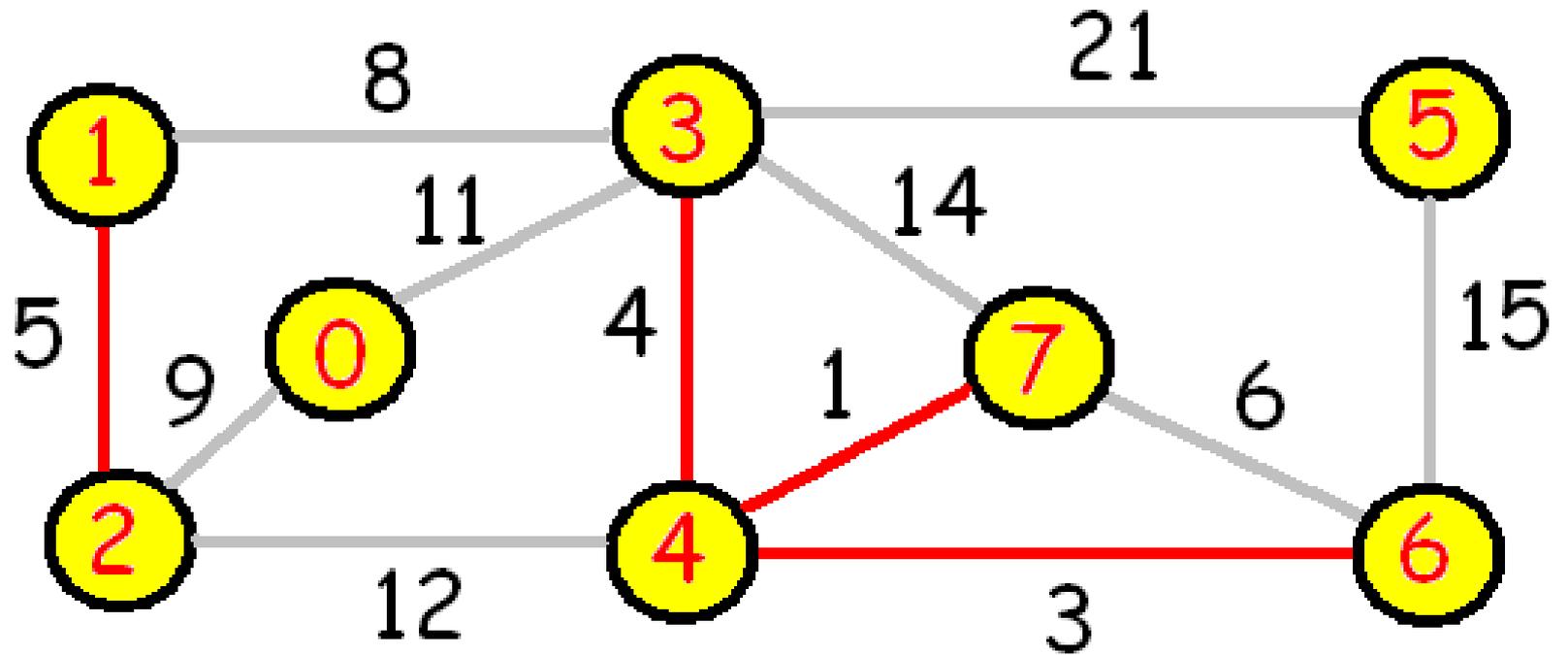
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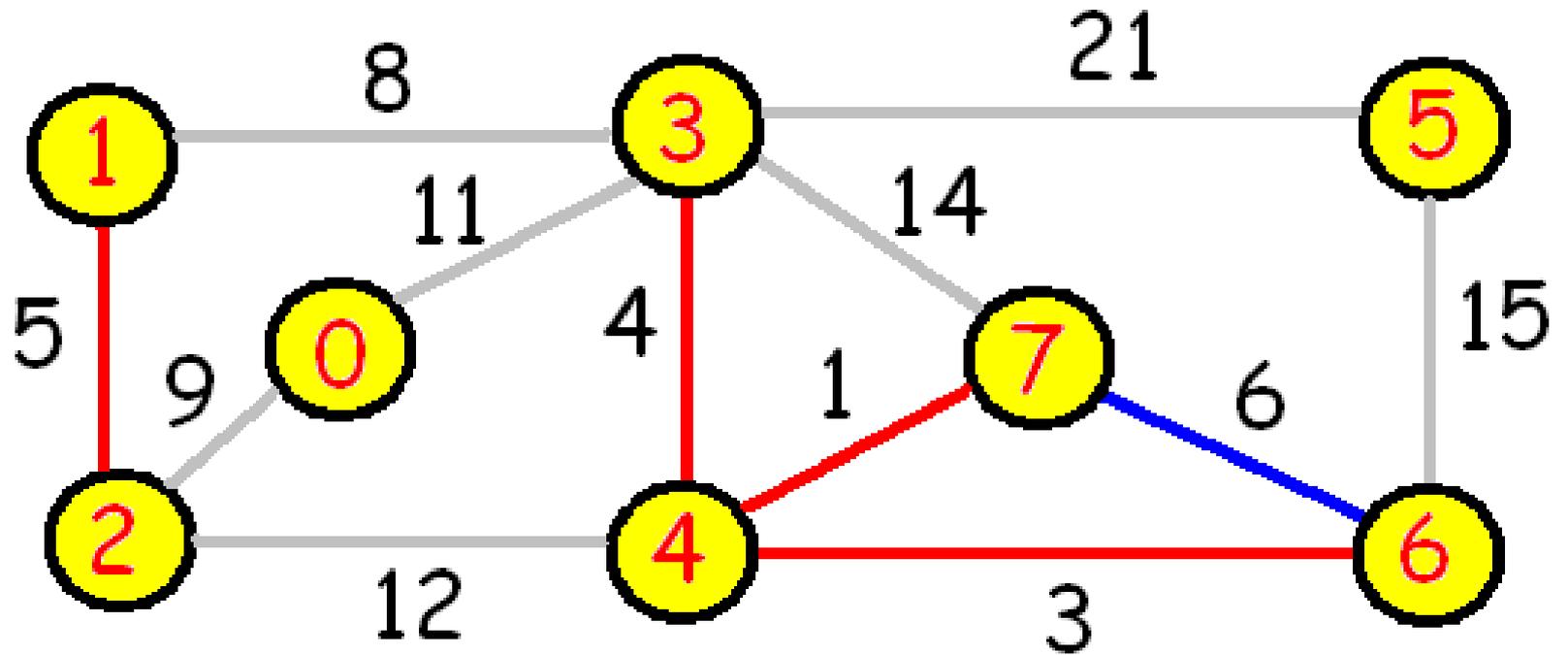
Edge weights:

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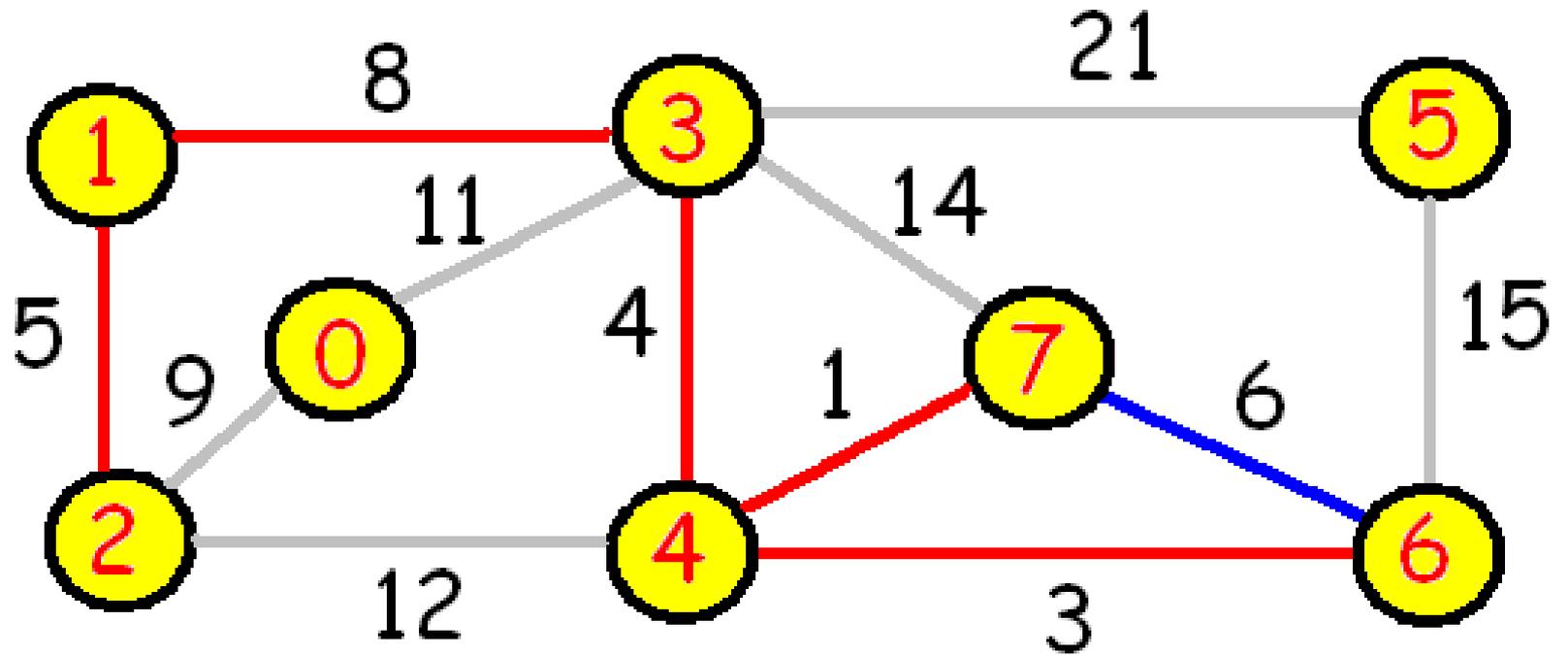
Edge weights:

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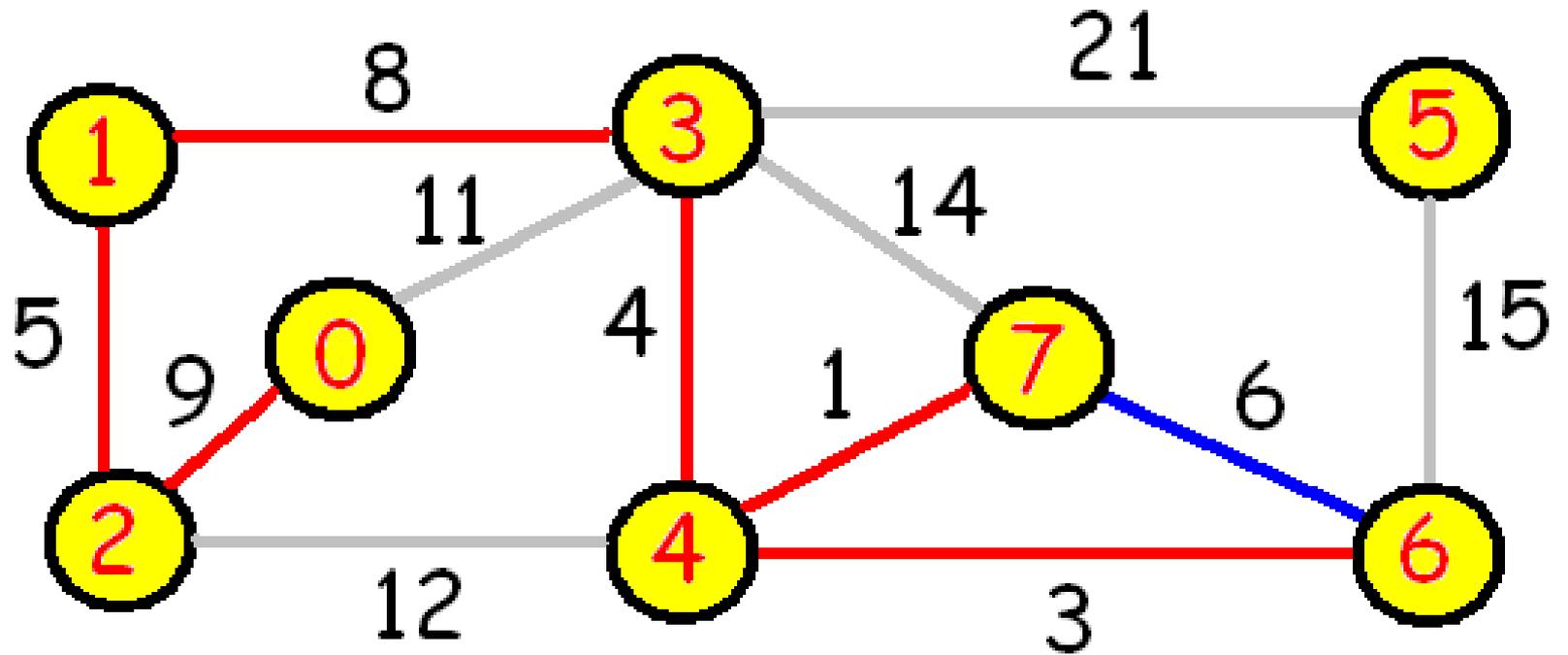
Edge weights:

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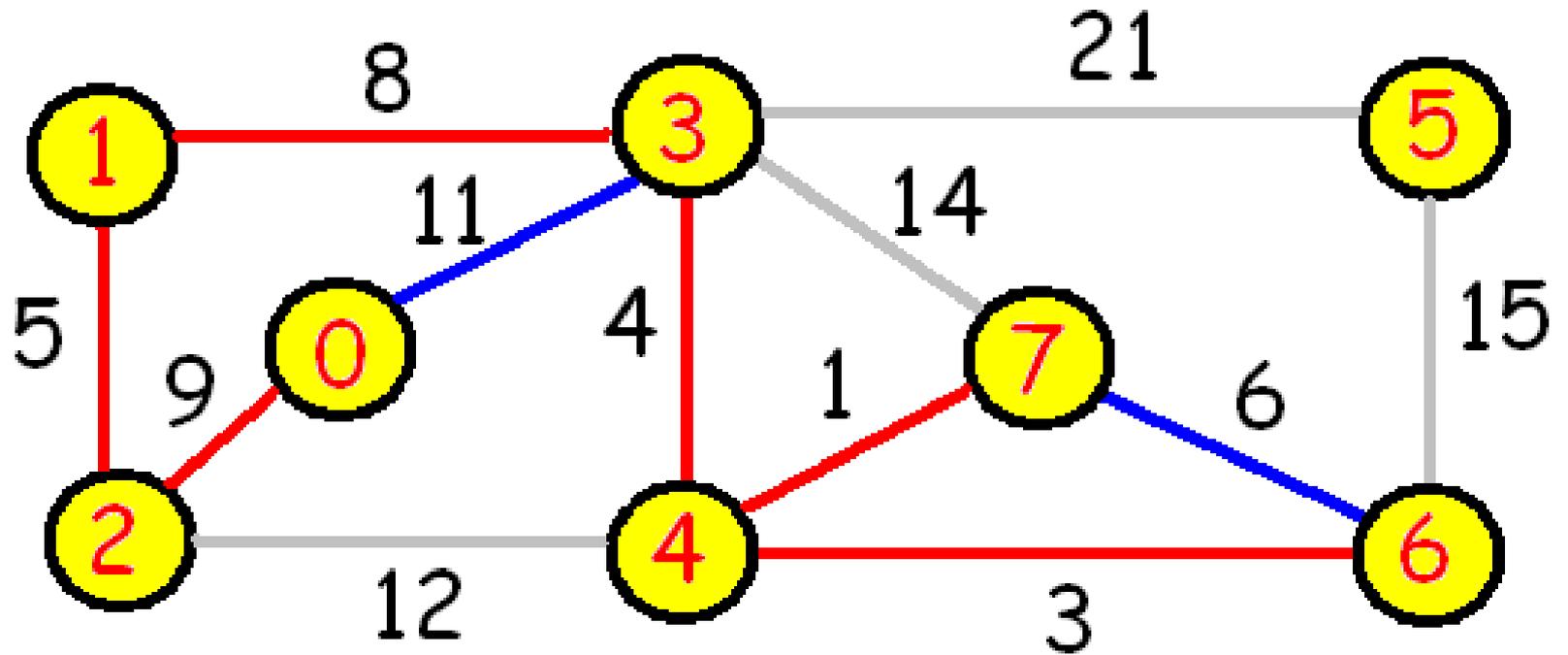
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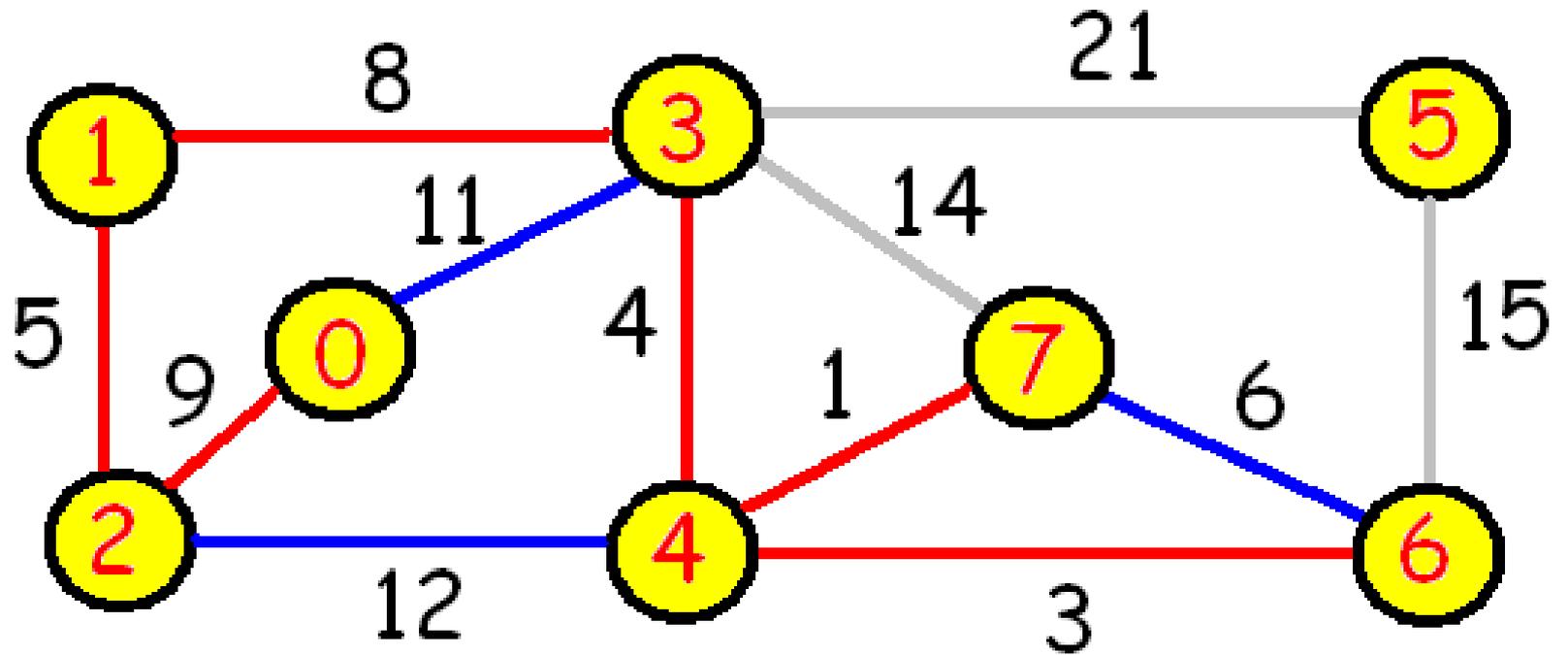
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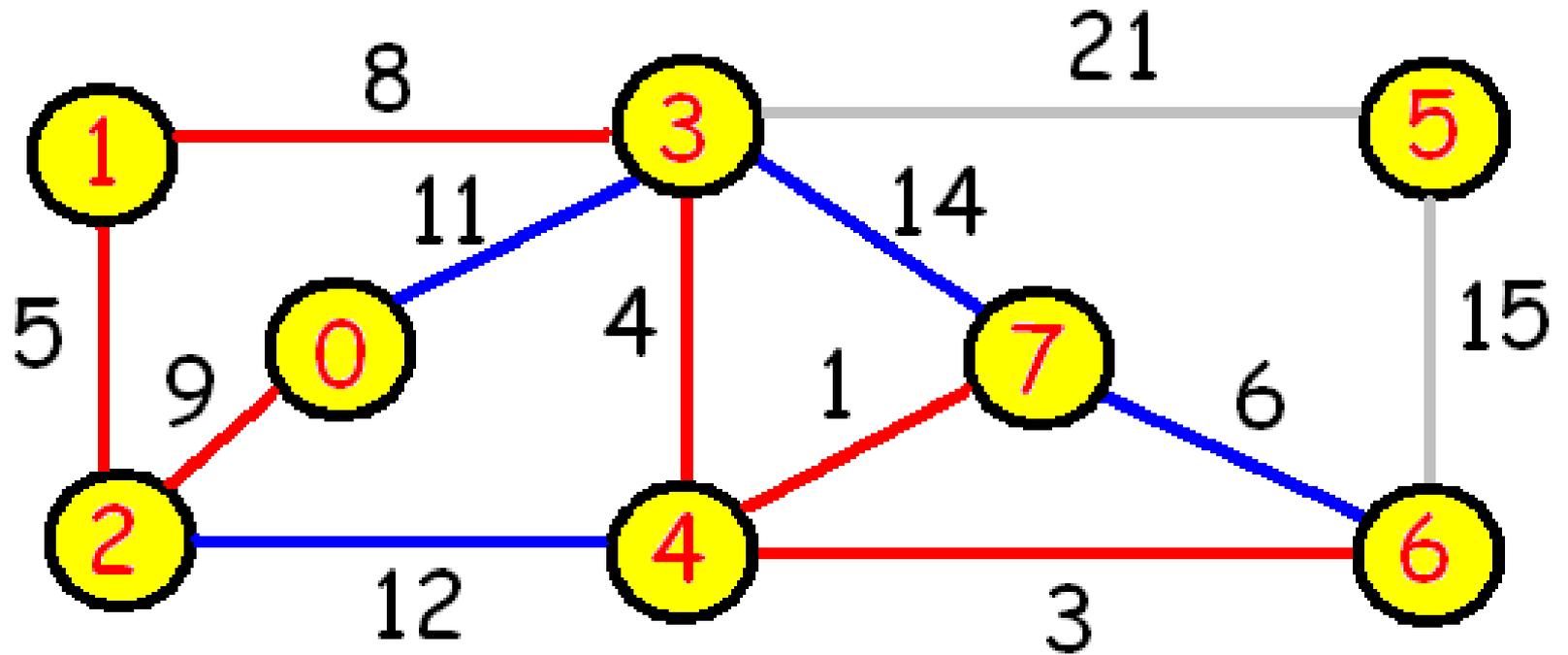
Edge weights:

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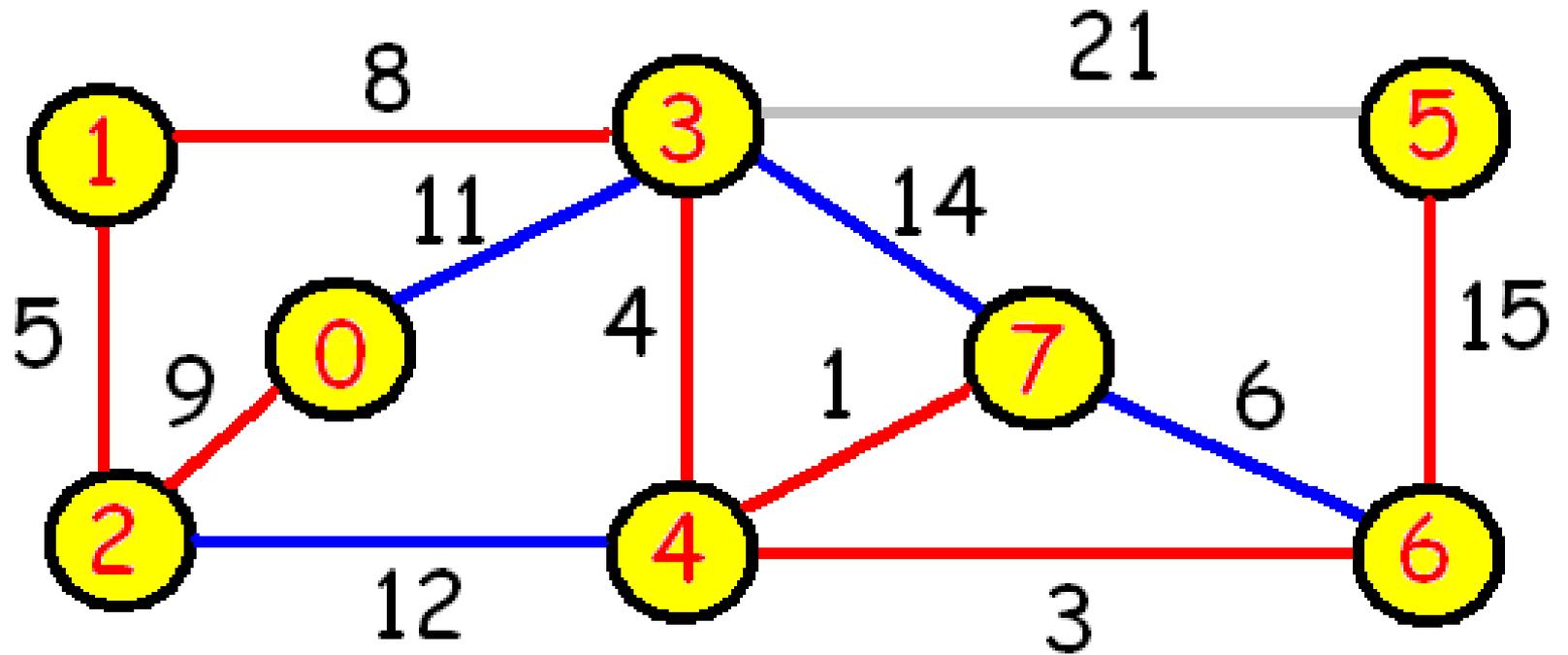
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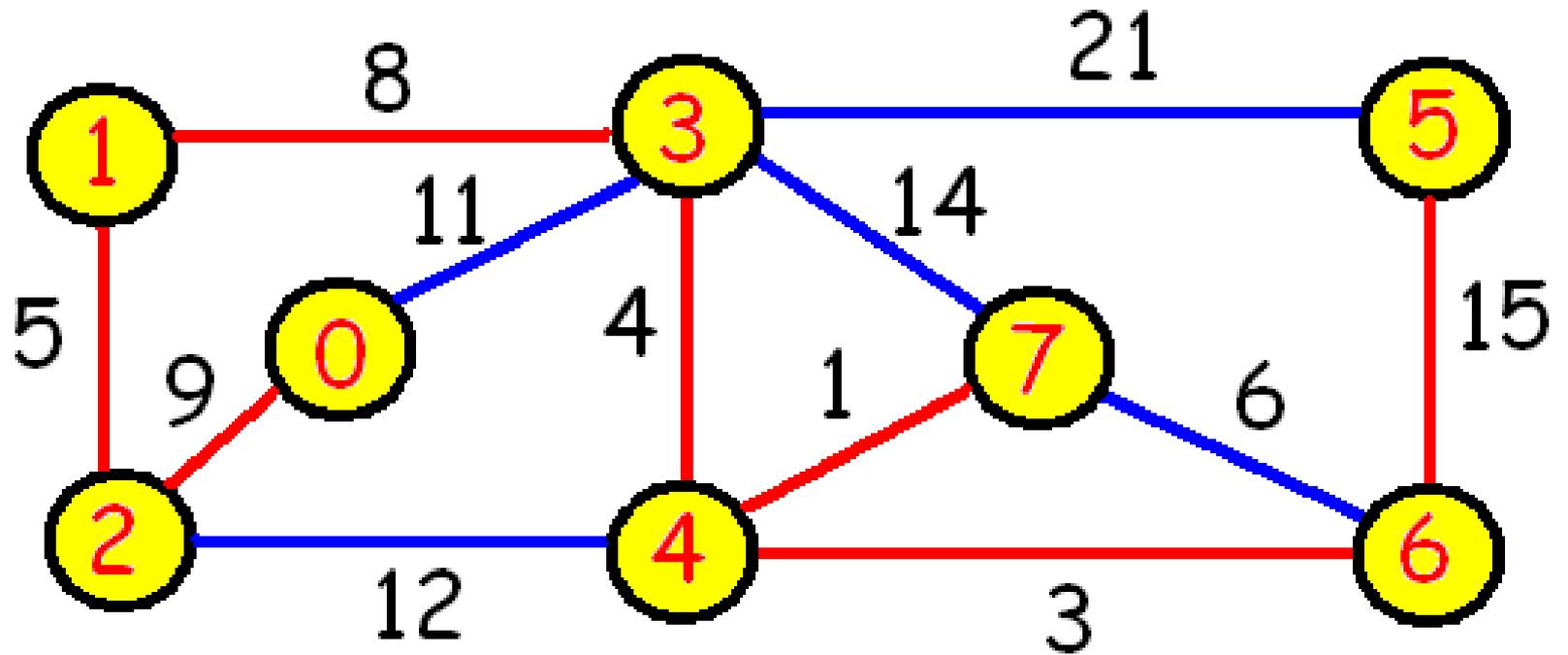
Edge weights:

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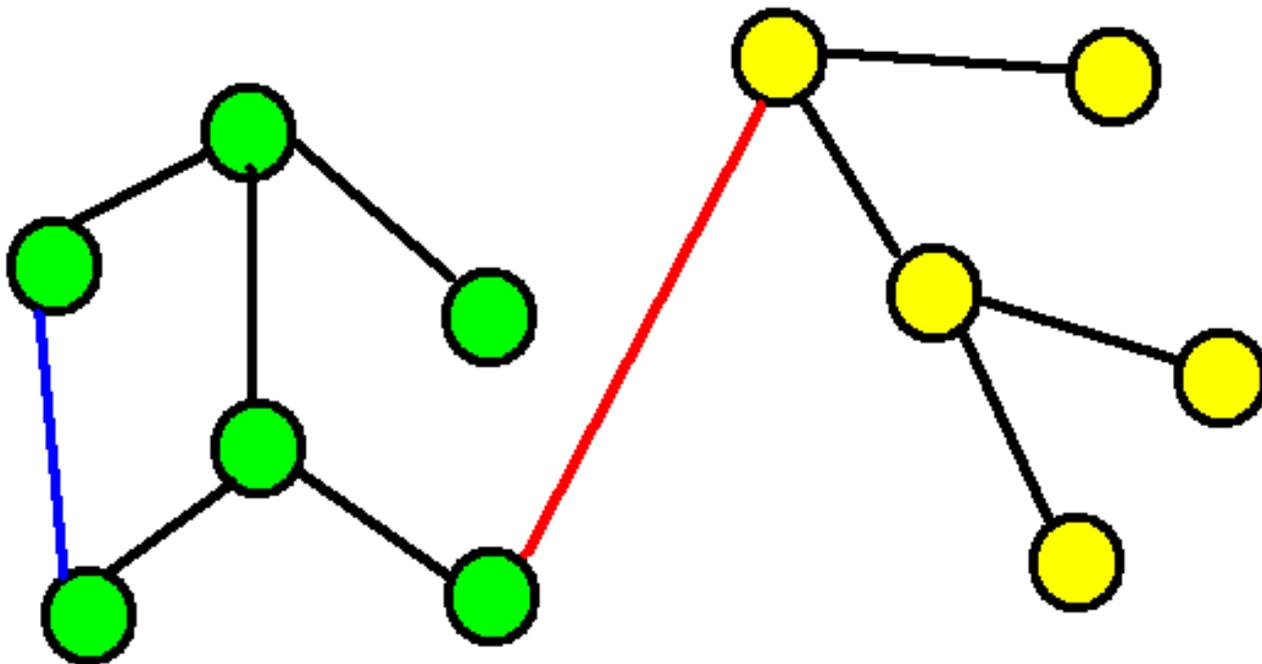


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How can we determine quickly at each step whether adding a new edge creates a cycle?

Or equivalently, given an edge  $(u,v)$  are  $u$  and  $v$  in the same component?



The UNION/FIND data structure is a dynamic data structure for graphs used to keep track of the connected components.

It has 2 routines:

**FIND(u)**: returns the name of the component containing vertex  $u$

**UNION(u, v)**: unions together the components containing  $u$  and  $v$  (corresponding to an addition of edge  $(u,v)$  to the graph).